A Large Term Rewrite System Modelling a Pioneering Cryptographic Algorithm

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Outline

- 1. Introduction to Term Rewrite Systems (TRS)
- 2. The Message Authenticator Algorithm (MAA)
- 3. Earlier Models of the MAA
- 4. Formal Modelling of the MAA as a TRS
- 5. Validation of the MAA Model
- 6. Conclusion



1. Introduction to Term Rewrite Systems (TRS)



Term Rewrite Systems (TRS)

A fundamental means to express computation

Basic concepts:

- sorts: abstract data domains
- operations: take N arguments and return one result
- terms: algebraic expressions (operations, free variables)
- ▶ rewrite rules: *left-hand term* \rightarrow *right-hand term* not (and (A, B)) \rightarrow or (not (A), not (B))
- Used in specification/programming languages
 - algebraic: abstract data types
 - functional: constructor types and pattern matching



Where can one find TRS models?

Paradox:

abundant literature on the theory of TRS

but difficult to find TRS models of realistic problems
 Available TRS models:

Rewrite Engines Contests (2006, 2008, 2010) the largest models have at most 300 lines

Specification of languages / compilers using TRS models can be large (10,000+ lines) but they are not "pure" TRS (they use strategies, sub-sorts, etc.)

This talk: a large TRS modelling a cryptographic algorithm



The REC Language

- REC: a textual notation for TRS models
- Introduced during the 2nd REC contest (2008)
 - human-readable, tool-independent format
 - supports strong typing (many-sorted specifications)
 - supports conditional rewrite rules (Boolean guards)
- We use a slightly enhanced version of REC
 - added distinction: constructors vs non-constructors
 - a few restrictions: left-linear rules, no equations between constructors, etc.
 - automatically translated into 13 different languages



Example 1: Booleans in REC

SORTS
Bool <
CONS
false : -> Bool < primitive operations
true : -> Bool (constructors)
OPNS
andBool : Bool Bool -> Bool < defined operations
orBool : Bool Bool -> Bool (non-constructors)
VARS
L : Bool < free variables
RULES
andBool (false, L) -> false $<$ rewrite rules that
andBool (true, L) -> L define non-constructors
orBool (false, L) -> L
orBool (true, L) -> true

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Example 2: Naturals in REC (1/2)

SORTS Nat CONS zero : -> Nat succ : Nat -> Nat OPNS addNat : Nat Nat -> Nat multNat : Nat Nat -> Nat eqNat : Nat Nat -> Bool ltNat : Nat Nat -> Bool VARS N N' : Nat



Example 2: Naturals in REC (2/2)

```
RULES
```

```
addNat (N, zero) -> N
addNat (N, succ (N')) -> addNat (succ (N), N')
```

```
multNat (N, zero) -> zero
multNat (N, succ (N')) -> addNat (N, multNat (N, N'))
```

```
eqNat (zero, zero) -> true
eqNat (zero, succ (N')) -> false
eqNat (succ (N), zero) -> false
eqNat (succ (N), succ (N')) -> eqNat (N, N')
```

ltNat (zero, zero) -> false
ltNat (zero, succ (N')) -> true
ltNat (succ (N'), zero) -> false
ltNat (succ (N), succ (N')) -> ltNat (N, N')



2. The Message Authenticator Algorithm (MAA)



Cryptography basics

Message Digest

- ▶ function: (long) message \rightarrow (short) numeric value
- ensures integrity (the message has not been modified)
- ▶ example: MD5

Message Authentication Code (MAC)

- ▶ function: (long) message, (short) key \rightarrow (short) value
- the key is secret, shared by the sender and the receiver
- ensures both authentication and integrity
- examples: hash-based (HMAC), universal (UMAC), block ciphers (CMAC, OMAC, PMAC), etc.



Message Authenticator Algorithm (MAA)

- First widely-used MAC function
- Designed by Donald Davies and David Clayden (NPL, 1983)



- to protect banking transactions
- intended to be implemented in software (32-bit PCs)
- Adopted by financial institutions
 - standardized by ISO in 1987 [ISO 8730 and 8731-2]
 - attacks published in the mid 90s
 - withdrawn from ISO standards in 2002

Overview of the MAA

Inputs:

- A 64-bit key (split into two blocks J, K)
- A message, seen as a sequence of blocks (message should be less than 1,000,000 blocks)

Outputs:

- A 32-bit MAC value (much too short nowadays!)
- Basic operations:
 - Iogical: AND, OR, XOR, CYC (bit rotation)
 - arithmetic: ADD, MUL (mod 2³²), MUL1 (mod 2³²-1), MUL2 (mod 2³²-2), MUL2A (faster variant of MUL2)



MAA data flow

Prelude: converts key (J, K) into 6 blocks X0, Y0, V0, W, S, T

Main Loop: iterates on each message block, modifying 3 variables X, Y, V

Coda: two final iterations on the two blocks S and T



"Mode of operation"

Message is split into a list of 256-block segments



3. Earlier Models of the MAA



Why choosing the MAA?

More challenging than conventional examples:

- protocols deal with simple data types
- compilers deal with abstract syntax trees (explored using standard traversals)
- cryptographic functions exhibit "strange" behavior by performing "irregular" calculations
- Large example, still of manageable complexity
- Definition of MAA is stable and available
- MAA played a role in the history of formal methods



Informal specifications

[Davies-Clayden-88] NPL technical report

- complete definition of the MAA
- gives two implementations in C and BASIC
- these implementation do not support "mode of operation" (only work for messages <= 256 blocks)
 [ISO standard 8731-2]
 - core part very similar to [Davies-Clayden-88]
- These definitions in natural language are ambiguous at several places
 - e.g. byte ordering, mode of operation



Formal specifications (1/2)

NPL chose MAA to assess formal methods

they developed 3 formal specifications of the MAA

1) VDM [G. I. Parkin and G. O'Neill, 1990]

- included as Annex B of ISO standard 8731-2:1992
- 3 implementations derived manually from VDM:
 C, Miranda, Modula-2

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- **2)** Z [M. K. F. Lai, 1991]
 - Knuth's "literate programming" approach
 - Z code fragments inserted in natural-language ISO text

Formal specifications (2/2)

3) LOTOS abstract data types [H. Munster, 1991]

- fully formal, but non executable
- "wishful thinking" equations: "given x, the result is y such that x = f (y)" ⇒ requires to invert function f
- 4) LOTOS abstract data types [H. Garavel, Ph. Turlier, 1992]
 - derived from [Munster-91]
 - rewritten to remove "wishful thinking" equations
 - a few types and functions implemented directly in C
 - implementation automatically derived (CAESAR.ADT)

Goals of our work

Provide a model of the MAA in REC language with (at least) five qualities:

Formal
Exhaustive
Self-contained
Validated
Executable

(no natural language)
(the full MAA is described)
(no external C code)
(correctness properties)
(implementations generated automatically in 13 languages)



4. Formal Modelling of the MAA as a TRS



Starting point

Informal description of the MAA

- [Davies-Clayden-88] NPL research report quasi identical as [ISO standard 8731-2]
- together with its C implementation although incomplete (no "mode of operation")

Formal description of the MAA

- [Garavel-Turlier-92] specification in LOTOS and C
- derived from the LOTOS specification of [Munster-91]



Outcome

- Formal model of the MAA as a TRS in REC language
- A large model:
 - 46 pages of text (Annex B of our paper)
 - 1575 lines (5 times larger than the largest benchmarks of the Rewrite Engines Competition)
 - ▶ 13 sorts
 - 18 constructors
 - 644 non-constructors
 - 684 rewrite rules

(only 6 conditional rules that can be easily eliminated)



Good properties

- Our model is exhaustive
 - it describes the full MAA (including "mode of operation")
- Our model is minimal
 - each sort, constructor, and non-constructor defined is actually used (no "dead code")
- Our model is self-contained
 - each detail of the MAA is expressed using TRS only
 - no import of externally-defined types or functions
 - no machine-specific assumptions (e.g., 32-bit vs 64-bit words, big-endian ordering)



Test vectors

Cryptographic functions come with test vectors

Our model is self-checking
it contains 202 accortions to

J	00FF	00FF	00FF	00FF	5555	5555	5555	5555
к	0000	0000	0000	0000	5 A3 5	D667	5A35	D667
Р		FF		FF		00		00
X ₀	4A64	5A01	4A64	5A01	34AC	F886	34AC	F886
Y	50DE	C930	50DE	C930	7397	C9AE	7397	C9AE
v _o	5CCA	3239	5CCA	3239	7201	F4DC	7201	F4DC
Ŵ	FECC	AA6E	FECC	AA6E	2829	040B	2829	040B
м,	5555	5555	ΑΑΑΑ	AAAA	0000	0000	FFFF	FFFF
x	48B2	04D6	6AEB	ACF8	2FD7	6FFB	8DC8	BBDE
Y	5834	A585	9DB1	5CF6	550D	91CE	FE4E	5BDD
M ₂	AAAA	AAAA	5555	5555	FFFF	FFFF	0000	0000
x	4F99	8E01	270E	EDAF	A70F	C148	CBC8	65BA
Y	BE9F	0917	B814	2629	1D10	D8D3	0297	AF6F
S	51ED	E9C7	51ED	E9C7	9E2E	7B36	9E2E	7B36
x	3449	25FC	2990	7CD8	B1CC	1CC5	3CF3	A7D2
Y	DB91	02B0	BA92	DB12	29C1	485F	160E	E9B5
т	24B6	6FB5	24B6	6FB5	1364	7149	1364	7149
x	277B	4B25	28EA	D8B3	288F	C786	D048	2465
Y	D636	250D	81D1	0CA3	9115	A558	7050	EC5E
Z	F14D	6E28	A93B	D410	B99A	62DE	A018	C83B

- it contains 203 assertions test vectors
 - taken from [Davies-Clayden-88], i.e., [ISO 8731-2]
 - taken from [ISO 8730:1990, Annex E.3.3]
 - added by us, so as to detect:
 - errors arising from byte permutations (endianness issues)
 - incorrect segmentation of messages longer than 256 blocks

Executability issues

- In principle, TRS encoded in the REC format are executable (by translation to other languages)
- In practice, Peano-style naturals (i.e., in unary notation with zero and succ) exhaust memory
 - the MAA manipulates many blocks (32-bit naturals)
 - blocks cannot be represented in unary notation
 - we represent blocks in binary form (words of 4 octets)
 - logical operations (AND, OR, XOR, CYC) are easy
 - ▶ arithmetical operations (ADD, CAR, MUL) are involved \Rightarrow 8-bit, 16-bit, and 32-bit adders and multipliers



Readability

- Our model is readable (despite its size)
 - regular naming conventions for all identifiers
 - constructors chosen appropriately
 - definitions of non-constructors kept simple
- Modular structure:
 - in the MARS repository: the MAA model is a monolithic REC file
 - in Annex B of our paper: the MAA model is split into 21 sections



Guided tour of the MAA model (1/3)

- 21 sections in Annex B of our paper
- BASIC SORTS
 - ▶ 1. Bool sort
 - ▶ 2. Nat sort (only used for "small" numbers \leq 4100)
- MACHINE WORDS
 - ▶ 3. Bit sort
 - ▶ 4. Octet sort (8 bits)
 - 5. OctetSum sort (9 bits: an Octet and a carry bit)
 - ▶ 6. Half sort (16 bits)



Guided tour of the MAA model (2/3)

- 7. HalfSum sort (17 bits: a Half and a carry bit)
- ▶ 8. Block sort (32 bits)
- 9. BlockSum sort (33 bits: a Block and a carry bit)
- 10. Pair sort (64 bits)
- INPUT/OUTPUT DATA
 - ▶ 11. Key sort (64 bits)
 - 12. Message sort (non-empty list of Blocks)
 - 13. SegmentedMessage sort (non-empty list of Messages, each containing at most 256 blocks)



Guided tour of the MAA model (3/3)

CRYPTOGRAPHIC FUNCTIONS

- ▶ 14. functions CYC, FIX1, FIX2, adjust, PAT, BYT, ADDC
- ▶ 15. functions MUL1, MUL2, MUL2A
- ▶ 16. functions Hi, J1_i, J2_i, K1_i, K2_i
- ► 17. Prelude, MainLoop, Coda, Segmentation

TEST VECTORS

- 18. Tables 1, 2, and 3 of [Davies-Clayden-88]
- 19. Table 4 of [Davies-Clayden-88] and other tests
- 20. Table 5 of [Davies-Clayden-88]
- > 21. Table 6 of [Davies-Clayden-88] and other tests

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5. Validation of the MAA Model



Properties

None of the prior formal MAA specifications (in VDM, Z, and LOTOS) was proven correct

Our REC specification brings stronger guarantees:

- confluence
- termination
- ► confluence and termination ⇒ all rewrite strategies produce the same result
- functional correctness of the 203 test vectors



Confluence and Termination

- Our TRS is deterministic, thus confluent
 - all constructors are free
 - all the rewrite rules that define a non-constructor have disjoint patterns and mutually exclusive premises
 - this was checked by the Opal compiler after automatic translation of the REC model into the Opal language.
- Our TRS is terminating
 - the REC model was automatically translated into the TRS input format of the AProVE tool
 - AProVE produced a proof of quasi-decreasingness (76 steps, 420 pages)

Functional correctness

- Our REC model was automatically translated into 13 languages: Clean, Haskell, LNT, LOTOS, Maude, mCRL2, OCaml, Opal, Rascal, Scala, Standard ML, Stratego/XT, and Tom
- It was then submitted to 16 tools (compilers, interpreters, and rewrite engines):
 - ▶ 11 tools reported that all the 203 test vectors pass
 - (the other tools gave up or timed out)
 - moreover, binary adders and multipliers have been checked separately using 30,000 test vectors



Two errors detected

Incorrect test vectors given for function PAT [Davies-Clayden-88, Table 3] and [ISO 8732-2:1992, Table A.3]

{X0,Y0}	0103	0703	1D3B	7760	PAT{X0,Y0}	EE
{VO,W}	0103	050B	1706	5DBB	PAT{VO,W}	BB
{S,T}	0103	0705	8039	7302	PAT{S,T}	E6

should read:

{H4,H5}	0000	0003	0000	0060	PAT{H4,H5}	EE
{H6,H7}	0003	0000	0006	0000	PAT{H6,H7}	BB
{H8,H9}	0000	0005	8000	0002	PAT{H8,H9}	E6

 Error in the handwritten C function provided to implement the LOTOS function HIGH_MUL
 ⇒ mixing formal and non-formal code is risky



6. Conclusion



Contributions

We revisited the Message Authenticator Algorithm

- an pioneering algorithm in cryptography (80s)
- an early application of formal methods (90s)
- We enriched the MARS model repository
 - a formal model of the MAA in the REC language
 - one of the largest handwritten TRS available today
 - self-contained and minimal
 - validated (confluence, termination, test vectors)
 - executable: translations into 13 different languages
 - reusable components (binary adders and multipliers)



Future work

Caution! our MAA model is a "tour de force"

- TRS do not scale well to large problems
- considerable effort was needed to produce a structured, readable REC model
- 2-6 times longer than any other (formal or informal) description of the MAA
- Possible uses of our MAA model
 - Iab exercises for students (see Annex B.22)
 - assessment of tools (e.g., 1÷140 speed ratio)
 - provers: verify correctness of binary adders/multipliers

