# Compositional verification applied to RERS 2019

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RERS 2019

## Approach

- Categories: parallel CTL and parallel LTL
- Main tool: CADP (cadp.inria.fr)
- Auxiliary tools:
  - SPOT (spot.lrde.epita.fr)
    Translation of LTL to Büchi automata
  - KandISTI/FMC (fmt.isit.cnr.it/kandisti)
    Cross-checking of CTL results
  - nuXmv (nuxmv.fbk.eu)
    Cross-checking of LTL results
- Main technique: Compositional verification

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#### The CADP toolbox http://cadp.inria.fr



- Developed by Inria/CONVECS for > 30 years
- Model & equivalence checking, rapid prototyping, test case generation, ... (> 80 tools and libraries)
- Enumerative techniques: LTS model
- Main languages and tools used in this work:
  - LNT system description language,
  - MCL property description language,
  - EVALUATOR model checker ,
  - BCG\_MIN LTS minimization tool,
  - SVL scripting language and compiler, …

#### **RERS parallel verification tasks**

- System description  $P_1 \mid | \dots | | P_n$ 
  - 9 system descriptions from 8 to 70 parallel processes and from 29 to 234 actions
  - We used the DOT representation
  - Automated translation from DOT to LNT
- Property  $\varphi$ 
  - 20 CTL properties for each system description
  - 20 LTL properties for each system description

#### **CTL compositional verification**

• Results of [MW14] are used to infer from  $\varphi$ 

– a set of actions *H* that can be hidden

– an equiv. relation R that preserves  $\varphi$  (improved)

- A reduced model *M* is obtained using SVL as smart *R* reduction of hide *H* in *P*<sub>1</sub> || ... || *P*<sub>n</sub>
- $\varphi$  is verified on *M* using EVALUATOR:  $P_1 \mid \mid \dots \mid \mid P_n \mid = \varphi$  iff  $M \mid = \varphi$

[MW14] R. Mateescu, A. Wijs. *Property-Dependent Reductions* Adequate With Divergence-Sensitive Branching Bisimilarity. SCP, 2014.

### **CTL results**

- All 180 CTL properties verified on this laptop:
  - 158 min. CPU (≈ 2.5 hours) / ≈ 5 hours elapsed
  - 200 MB memory
  - Largest intermediate LTS  $\leq$  3363 states
- Cross-checking with KandISTI/FMC:
  - on the fly, explicit verification on unreduced LTS
  - 126 problems solved out of 180 (70 %)
    max 2h, 64 GB memory available
  - CADP results confirmed

#### LTL compositional verification

- Reduced model *M* obtained using same approach
- Use of Büchi automaton **B** 
  - Automated translation of  $\neg \varphi$  to transition-based Büchi automaton using SPOT (HOA format)
  - Automated encoding from HOA to LNT
  - Accepting transitions encoded by action ACC
- EVALUATOR is used to verify the acceptance condition encoded as an MCL formula:

 $P_1 || ... || P_n |= \varphi$  iff  $M || B |= \neg < true^*$ . ACC> @

### LTL results

- All 180 LTL properties verified on this laptop:
  - 144 min. CPU (≈ 2.5 hours) / ≈ 5 hours elapsed
  - 200 MB memory
  - Largest intermediate LTS  $\leq$  1068 states
- Cross-checking with nuXmv:
  - LTL verification <u>on the reduced LTS</u> (risk)
  - all problems solved
  - CADP results confirmed

#### Conclusion

- Compositional verification is effective to solve CTL and LTL parallel benchmarks of RERS 2019
- Causes of success:
  - Expressive languages (LNT, MCL, SVL, ...)
  - Efficient tools
  - Team working: combination of expertises, synergy
  - Hard work and tenacity
- Diversity of approaches  $\Rightarrow$  trust increases
- New results : papers in preparation